

First Order Logic: Conversion to CNF

1. Eliminate biconditionals and implications:

- Eliminate \Leftrightarrow , replacing $\alpha \Leftrightarrow \beta$ with $(\alpha \Rightarrow \beta) \wedge (\beta \Rightarrow \alpha)$.
- Eliminate \Rightarrow , replacing $\alpha \Rightarrow \beta$ with $\neg\alpha \vee \beta$.

2. Move \neg inwards:

- $\neg(\forall x p) \equiv \exists x \neg p$,
- $\neg(\exists x p) \equiv \forall x \neg p$,
- $\neg(\alpha \vee \beta) \equiv \neg\alpha \wedge \neg\beta$,
- $\neg(\alpha \wedge \beta) \equiv \neg\alpha \vee \neg\beta$,
- $\neg\neg\alpha \equiv \alpha$.

3. Standardize variables apart by renaming them: each quantifier should use a different variable.

4. Skolemize: each existential variable is replaced by a *Skolem constant* or *Skolem function* of the enclosing universally quantified variables.

- For instance, $\exists x Rich(x)$ becomes $Rich(G1)$ where $G1$ is a new Skolem constant.
- “Everyone has a heart” $\forall x Person(x) \Rightarrow \exists y Heart(y) \wedge Has(x, y)$
becomes $\forall x Person(x) \Rightarrow Heart(H(x)) \wedge Has(x, H(x))$,
where H is a new symbol (Skolem function).

5. Drop universal quantifiers

- For instance, $\forall x Person(x)$ becomes $Person(x)$.

6. Distribute \wedge over \vee :

- $(\alpha \wedge \beta) \vee \gamma \equiv (\alpha \vee \gamma) \wedge (\beta \vee \gamma)$.